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| | 7590 02/27/2007 FEPHENSON ASCOLE | EXAMINER | | |
| 4807 SPICEWOOD SPRINGS RD. BLDG. 4, SUITE 201 AUSTIN, TX 78759 | | | COLIN, CARL G | |
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| SHORTENED STATUTORY | Y PERIOD OF RESPONSE | MAIL DATE | DELIVERY MODE | |
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Please find below and/or attached an Office communication concerning this application or proceeding.

If NO period for reply is specified above, the maximum statutory period will apply and will expire 6 MONTHS from the mailing date of this communication.

| | Application No. | Applicant(s) | | | | |
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| | 09/456,692 | FRY ET AL. | | | | |
| Office Action Summary | Examiner | Art Unit | | | | |
| • | Carl Colin | 2136 | | | | |
| The MAILING DATE of this communication appears on the cover sheet with the correspondence address Period for Reply | | | | | | |
| A SHORTENED STATUTORY PERIOD FOR REPLY WHICHEVER IS LONGER, FROM THE MAILING DA - Extensions of time may be available under the provisions of 37 CFR 1.13 after SIX (6) MONTHS from the mailing date of this communication. - If NO period for reply is specified above, the maximum statutory period w - Failure to reply within the set or extended period for reply will, by statute, Any reply received by the Office later than three months after the mailing earned patent term adjustment. See 37 CFR 1.704(b). | ATE OF THIS COMMUNICATION B6(a). In no event, however, may a reply be tir rill apply and will expire SIX (6) MONTHS from cause the application to become ABANDONE | N. nely filed the mailing date of this communication. D (35 U.S.C. § 133). | | | | |
| Status | | | | | | |
| 1) Responsive to communication(s) filed on <u>01 De</u> | ecember 2006. | | | | | |
| ·— | ,— | | | | | |
| • • | 3) Since this application is in condition for allowance except for formal matters, prosecution as to the merits is | | | | | |
| closed in accordance with the practice under Ex parte Quayle, 1935 C.D. 11, 453 O.G. 213. | | | | | | |
| Disposition of Claims | | | | | | |
| 4) Claim(s) 107-147 and 165-181 is/are pending i 4a) Of the above claim(s) is/are withdray 5) Claim(s) is/are allowed. 6) Claim(s) 107-147 and 165-181 is/are rejected. 7) Claim(s) is/are objected to. 8) Claim(s) are subject to restriction and/or | vn from consideration. | | | | | |
| Application Papers | | | | | | |
| 9) The specification is objected to by the Examine 10) The drawing(s) filed on is/are: a) acce Applicant may not request that any objection to the Replacement drawing sheet(s) including the correct 11) The oath or declaration is objected to by the Ex | epted or b) objected to by the drawing(s) be held in abeyance. Se ion is required if the drawing(s) is ob | e 37 CFR 1.85(a). ejected to. See 37 CFR 1.121(d). | | | | |
| Priority under 35 U.S.C. § 119 | | | | | | |
| 12) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f). a) All b) Some * c) None of: 1. Certified copies of the priority documents have been received. 2. Certified copies of the priority documents have been received in Application No. 3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)). * See the attached detailed Office action for a list of the certified copies not received. | | | | | | |
| Attachment(s) 1) Notice of References Cited (PTO-892) 2) Notice of Draftsperson's Patent Drawing Review (PTO-948) 3) Information Disclosure Statement(s) (PTO/SB/08) Paper No(s)/Mail Date | 4) Interview Summary Paper No(s)/Mail D 5) Notice of Informal I 6) Other: | Pate | | | | |

DETAILED ACTION

1. In view of the Appeal Brief filed on 12/1/2006, PROSECUTION IS HEREBY REOPENED. A new ground of rejection is set forth below.

To avoid abandonment of the application, appellant must exercise one of the following two options:

- (1) file a reply under 37 CFR 1.111 (if this Office action is non-final) or a reply under 37 CFR 1.113 (if this Office action is final); or,
- (2) initiate a new appeal by filing a notice of appeal under 37 CFR 41.31 followed by an appeal brief under 37 CFR 41.37. The previously paid notice of appeal fee and appeal brief fee can be applied to the new appeal. If, however, the appeal fees set forth in 37 CFR 41.20 have been increased since they were previously paid, then appellant must pay the difference between the increased fees and the amount previously paid.

A Supervisory Patent Examiner (SPE) has approved of reopening prosecution by signing below.

Response to Arguments

- 2. In response to communications filed on 12/1/2006, the following claims 107-147 and 165-181 are presented for examination.
- 2.1 In response to communications filed on 12/1/2006, the 112 rejection has been withdrawn.

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2.2 Applicant's arguments, filed in the appeal brief on 12/1/2006 have been fully considered but they are not moot in view of a new ground of rejection.

Claim Rejections - 35 USC § 103

- 3. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:
- (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

Claims 107-115, 131-139, and 165-173 are rejected under 35 U.S.C. 103(a) as being unpatentable over US Patent 6,064,671 to Killian in view of US Patent 6,643,701 to Aziz et al.

As per claim 107, Killian discloses a method comprising: providing a plurality of sockets in a socket table or array, wherein each socket normally represents a connection (i.e. has an associated connection) and each socket has an associated network port and network IP address (see column 14, lines 52-64 and column 15, lines 60-67 and column 5, lines 15-22) that meets the recitation of providing a plurality of sockets, wherein each socket has an associated connection and an associated security token (network address);

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Killian discloses the network address is provided by the associated connection (see column 6, lines 31-36 and lines 41-50) that meets the recitation of the associated security token is provided by the associated connection;

Killian also discloses receiving an outgoing message having a network address (see column 7, line 65 through column 8, line 2) and the message is associated with a connection (see column 29, lines 41-42), the message is also associated with a socket (see column 28, lines 26-27), that meets the recitation of *receiving a first connection and a first security token*;

Killian further discloses creating a socket and each socket has an associated connection (see column 14, lines 60-64 and column 5, lines 15-33) that meets the recitation of *creating a* socket associated with the first connection,

wherein the first connection has associated the first security token (network address) (see column 14, lines 60-64 and column 5, lines 15-33);

Killian further discloses comparing the network address with the associated network addresses for a match in the socket table (see column 16, lines 6-12) that meets the recitation of comparing the first security token with the associated security tokens;

in response to said comparing if none of the associated security tokens match the first security token (see column 16, lines 17-20), including the socket in the plurality of sockets (see column 16, lines 25-28).

Although **Killian's** disclosure referring to message is interpreted as a connection, it is obvious to one of ordinary skill in the art that the message is directed to a connection as explained above because it is associated with socket and as disclosed also in column 5, lines 52-60, the message is transmitted over a connection. It would have been obvious to one of ordinary

skill in the art at the time the invention was made to use the disclosure of Killian of routing message to route connection between physical networks for internetworking (see column 1, lines 45-65). Aziz et al in an analogous art teaches end to end security link by creating a first end-to-end security link between the first computer and a third computer and creating a second end-to-end security link between the second computer and the third computer to establish the secure connection (see column 3, lines 35-56 and column 7, line 56 through column 8, line 5).

Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to combine Killian and Aziz et al to create two separate connections at each end system to provide an application-specific authentication and link them to establish a secure connection between two end systems when messages are exchanged between computers.

As per claim 108, the references as combined above disclose the limitation of wherein the security token is one of a password, a network address, and a verification string (see Killian column 14, lines 52-64).

As per claim 109, Killian discloses in response to said comparing when a match occurs coupling an end point of the first connection to a network interface or to a machine which has a connection to the network at a remote location (see column 19, lines 4-7 and column 28, lines 38-46). Aziz et al discloses linking an end point of a first connection to an end point of a second connection in response to authentication (see column 7, line 49 through column 8, line 5), the authentication may include an authentication token such as password (see column 2, lines 10-15). Claim 109 is rejected on the same rationale as the rejection of claim 107.

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As per claim 110, the references as combined above disclose the limitation of in response to said comparing, if none of the associated security tokens match the first security token, upon a determination that the first connection is not to be associated with a socket, disconnecting the first connection (see Killian column 22, lines 43-47).

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As per claim 111, the references as combined above disclose wherein the coupling the first connection to the connection associated with the socket comprises: relaying a data stream between the first connection and the connection associated with the socket (see Killian, column 28, lines 38-46) and (see Aziz et al, column 7, line 49 through column 8, line 5).

As per claim 112, the references as combined above disclose wherein the coupling the first connection to the connection associated with the socket comprises: creating a single connection comprising the first connection and the connection associated with the socket (see Killian, column 19, lines 4-7 and column 28, lines 38-46) and (see Aziz et al, column 7, line 65 through column 8, line 1).

As per claim 113, the references as combined above disclose decoupling the first connection and the connection associated with the socket (see Killian, column 22, lines 30-42). As per claim 114, the references as combined above disclose the decoupling occurs upon one of failure and disconnect of one of the first connection and the connection associated with the socket (see Killian, column 22, lines 42-47).

As per claim 115, the references as combined above disclose wherein the first connection is transmitted through a first firewall program (see Killian, column 24, lines 48-59).

As per claim 131, Killian discloses an apparatus comprising: figures 8 and 13 for instance show computer systems for performing the claimed method of claim 1 comprising means for providing a plurality of sockets in a socket table or array, wherein each socket normally represents a connection (i.e. has an associated connection) and each socket has an associated network port and network IP address (see column 14, lines 52-64 and column 15, lines 60-67 and column 5, lines 15-22) that meets the recitation of means for providing a plurality of sockets, wherein each socket has an associated connection and an associated security token (network address);

Killian discloses the network address is provided by the associated connection (see column 6, lines 31-36 and lines 41-50) that meets the recitation of the associated security token is provided by the associated connection;

Killian also discloses end system and router (computer system) including means for receiving an outgoing message having a network address (see column 7, line 65 through column 8, line 2 and figs. 3-5) and the message is associated with a connection (see column 29, lines 41-

42), the message is also associated with a socket (see column 28, lines 26-27), that meets the recitation of means for receiving a first connection and a first security token;

Killian further discloses computer system with means for creating a socket and each socket has an associated connection (see column 14, lines 60-64 and column 5, lines 15-33 and fig. 3 and fig. 13) that meets the recitation of means for creating a socket associated with the first connection,

wherein the first connection has associated the first security token (network address) (see column 14, lines 60-64 and column 5, lines 15-33);

Killian further discloses means for comparing the network address with the associated network addresses for a match in the socket table (see column 16, lines 6-12 and fig. 3) that meets the recitation of means for comparing the first security token with the associated security tokens;

in response to said comparing if none of the associated security tokens match the first security token (see column 16, lines 17-20), means for including the socket in the plurality of sockets (see column 16, lines 25-28).

Although **Killian's** disclosure referring to message is interpreted as a connection, it is obvious to one of ordinary skill in the art that the message is directed to a connection as explained above because it is associated with socket and as disclosed also in column 5, lines 52-60, the message is transmitted over a connection. It would have been obvious to one of ordinary skill in the art at the time the invention was made to use the disclosure of Killian of routing message to route connection between physical networks for internetworking (see column 1, lines 45-65). **Aziz et al** in an analogous art teaches end to end security link by creating a first end-to-

end security link between the first computer and a third computer and creating a second end-toend security link between the second computer and the third computer to establish the secure
connection (see column 3, lines 35-56 and column 7, line 56 through column 8, line 5).

Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention
was made to combine Killian and Aziz et al to create two separate connections at each end
system to provide an application-specific authentication and link them to establish a secure
connection between two end systems when messages are exchanged between computers.

As per claim 132, the references as combined above disclose the limitation of wherein the security token is one of a password, a network address, and a verification string (see Killian column 14, lines 52-64).

As per claim 133, Killian discloses in response to said comparing when a match occurs coupling an end point of the first connection to a network interface or to a machine which has a connection to the network at a remote location (see column 19, lines 4-7 and column 28, lines 38-46). Aziz et al discloses a relay (220 and 440, figures 2 and 4) having means for linking an end point of a first connection to an end point of a second connection in response to authentication (see column 7, line 49 through column 8, line 5), the authentication may include an authentication token such as password (see column 2, lines 10-15). Claim 133 is rejected on the same rationale as the rejection of claim 131.

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As per claim 134, the references as combined above disclose the limitation of in response to said comparing, if none of the associated security tokens match the first security token, upon a determination that the first connection is not to be associated with a socket, means for disconnecting the first connection (see Killian column 22, lines 43-47).

As per claim 135, the references as combined above disclose wherein the coupling the first connection to the connection associated with the socket comprises: means for relaying a data stream between the first connection and the connection associated with the socket (see Killian, column 28, lines 38-46) and (see Aziz et al, column 7, line 49 through column 8, line 5).

As per claim 136, the references as combined above disclose wherein the coupling the first connection to the connection associated with the socket comprises: computer system such as the one shown in fig. 13 with means for creating a single connection comprising the first connection and the connection associated with the socket (see Killian, column 19, lines 4-7 and column 28, lines 38-46) and (see Aziz et al, column 7, line 65 through column 8, line 1).

As per claim 137, the references as combined above disclose means for decoupling the first connection and the connection associated with the socket (see Killian, column 22, lines 30-42 and fig. 13, 90B).

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As per claim 138, the references as combined above disclose the decoupling occurs upon one of failure and disconnect of one of the first connection and the connection associated with the socket (see Killian, column 22, lines 42-47).

As per claim 139, the references as combined above disclose wherein the first connection is transmitted through a first firewall program (see Killian, column 24, lines 48-59).

As per claim 165, claim 165 discloses the same limitations as claim 107 except for incorporating the claimed method into a computer program. Killian discloses the computer systems include applications for performing the invention (see column 14, lines 44-64).

Therefore, claim 165 is rejected on the same rationale as the rejection of claim 107.

As per claim 166, the references as combined above disclose the limitation of wherein the security token is one of a password, a network address, and a verification string (see Killian column 14, lines 52-64).

Claims 167-173 are similar to the rejected claims 109-115 respectively except for incorporating the claimed methods into a computer program. Therefore, 109-115 are rejected on the same rationale as the rejection of claims 167-173.

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4. Claims 116-119 are rejected under 35 U.S.C. 103(a) as being unpatentable over US Patent 6,064,671 to Killian in view of US Patent 6,643,701 to Aziz et al as applied to claims 107-115 and further in view of US Patent 6,104,716 to Crichton et al.

As per claims 116-119, Killian discloses a protocol daemon that can create network connections (column 21, lines 24-30 and lines 45-52) and discloses the first program providing a security token (see column 22, lines 8-20) but does not explicitly disclose the protocol daemon couples the first connection to the second connection. Aziz et al discloses a relay program for relaying first connection with second connection (see abstract and column 9, lines 30-39), either the client, the server, or the relay may be a protocol daemon (see column 6, lines 4-6), but does not explicitly disclose a second connection connecting the protocol daemon to the first program. Crichton et al in an analogous art teaches using a client proxy for communicating with a client and with a middle proxy and coupling the connections to provide end-to-end connections through firewalls (column 2, lines 26-52). Crichton et al also discloses the client and the proxy can reside on the same machine (column 6, lines 15-24). Crichton et al also discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). Crichton et al discloses one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). This means if the first program represents an application server an in-bound connection is created "a server end-proxy can connect to an inside X-Windows system server and a middle proxy" (column 5, lines 32-35). Applicant's specification discloses the same (on page 9, lines 9-15)

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program 135 (first program) requires an in-bound connection (e.g. where program 135 is an application server) ... such functionality is provided by a daemon running on computer 105. Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to modify the proxy or protocol daemon disclosed in **Killian** and **Aziz et al** as combined above to provide a protocol daemon program that does the creating of the first connection as well as the second connection and coupling the first and second connections thus increasing the functionality of end proxy to allow for other protocols and services as suggested by **Crichton et al** (see column 5, lines 32-45). One skilled in the art would have been lead to make such a modification and recognizes the advantage of using an end proxy that could provide both client and server end proxy functionality as this increase of functionality would allow for more protocols and services as suggested by **Crichton et al** (see column 5, lines 41-45).

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5. Claims 120-124, 128-130, 140-144, and 174-178 are rejected under 35 U.S.C. 103(a) as being unpatentable over US Patent 6,643,701 to Aziz et al in view of US Patent 6,104,716 to Crichton et al.

As per claim 120, Aziz et al discloses a method comprising: receiving from a client or first computer (first program) a first security token through handshaking (see column 5, lines 57-61 and column 8, lines 33-37) similar to the one described in fig. 1 (see column 1, line 55 through column 2, line 36) that meets the recitation of receiving a first security token from the first program;

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creating a second connection between relay and server (see column 8, lines 17-19 and column 10, lines 61-62) that meets the recitation of creating a second connection to a relay program;

providing the first security token to the relay program, for example (see column 5, lines 10-13 and column 9, lines 3-5); and upon successful creation of the second connection, coupling the first connection to the second connection, for example (see column 8, lines 50-65).

Aziz et al does not explicitly disclose creating a first connection to a first program. Crichton et al in an analogous art teaches using a client proxy for communicating with a client and with a middle proxy and coupling the connections to provide end-to-end connections through firewalls (column 2, lines 26-52). Crichton et al also discloses the client and the proxy can reside on the same machine (column 6, lines 15-24). Crichton et al also discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). Crichton et al discloses one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). This means if the first program represents an application server an in-bound connection is created "a server end-proxy can connect to an inside X-Windows system server and a middle proxy" (column 5, lines 32-35). Applicant's specification discloses the same (on page 9, lines 9-15) program 135 (first program) requires an in-bound connection (e.g. where program 135 is an application server) ... such functionality is provided by a daemon running on computer 105. Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to modify Aziz et al to provide a protocol daemon program that does the creating of

the first connection to a first program as well as the second connection and coupling the first and second connections thus increasing the functionality of end proxy to allow for other protocols and services as suggested by **Crichton et al** (see column 5, lines 32-45). One skilled in the art would have been lead to make such a modification and recognizes the advantage of using an end proxy that could provide both client and server end proxy functionality as this increase of functionality would allow for more protocols and services as suggested by **Crichton et al** (see column 5, lines 41-45).

As per claim 121, the references as combined above disclose the limitation of wherein the second connection is transmitted through a firewall program (see Crichton et al, fig. 4).

Claim 121 is rejected on the same rationale as the rejection of claim 120.

As per claim 122, the references as combined above disclose relaying a data stream between the first connection and the second connection (see Aziz et al, column 8, lines 50-65).

As per claim 123, the references as combined above disclose wherein the first security token is one of a password, a network address, and a verification string (see Aziz et al, column 9, lines 3-5).

As per claim 124, the references as combined above disclose terminating the first connection and the second connection (see Aziz et al, column 8, lines 48-49).

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above.

As per claims 128-130, Aziz et al discloses receiving a first security token from the first program (see column 5, lines 57-61 and column 8, lines 33-37; providing the first security token to the relay program, for example (see column 5, lines 10-13 and column 9, lines 3-5) a relay program for relaying first connection with second connection (see abstract and column 9, lines 30-39 and column 8, lines 50-65), and further discloses either the client, the server, or the relay may be a protocol daemon and may be incorporated in one or more machines (see column 6, lines 4-6). Aziz et al does not explicitly disclose a protocol daemon does the creating the first connection, the creating the second connection. Crichton et al in an analogous art teaches using a client proxy for communicating with a client and with a middle proxy and coupling the connections to provide end-to-end connections through firewalls (column 2, lines 26-52). Crichton et al also discloses the client and the proxy can reside on the same machine (column 6, lines 15-24). Crichton et al also discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). Crichton et al discloses the protocol daemon creating connection as explained in

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As per claim 140, Aziz et al discloses an apparatus comprising: a relay (220 and 440, figures 2 and 4) having means for receiving from a client or first computer (first program) a first security token through handshaking (see column 5, lines 57-61 and column 8, lines 33-37) similar to the one described in fig. 1 (see column 1, line 55 through column 2, line 36) that

claim 12 above. Claims 128-130 are rejected on the same rationale as the rejection of claim 120

meets the recitation of means for receiving a first security token from the first program (column 5, lines 18-20);

the relay (220 and 440, figures 2 and 4) having means for creating a second connection between relay and server (see column 8, lines 17-19 and column 10, lines 61-62) that meets the recitation of means for creating a second connection to a relay program;

means for providing the first security token to the relay program, for example (see column 5, lines 10-13 and column 9, lines 3-5); and upon successful creation of the second connection, means for coupling the first connection to the second connection, for example (see column 8, lines 50-65).

Aziz et al does disclose that the relay may be used as a daemon (column 6, lines 4-6) but is silent about means for creating a first connection to a first program, Crichton et al in an analogous art teaches using a client proxy for communicating with a client and with a middle proxy and coupling the connections to provide end-to-end connections through firewalls (column 2, lines 26-52). Crichton et al also discloses the client and the proxy can reside on the same machine (column 6, lines 15-24). Crichton et al also discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). Crichton et al discloses one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45). This means if the first program represents an application server an in-bound connection is created "a server end-proxy can connect to an inside X-Windows system server and a middle proxy" (column 5, lines 32-35). Applicant's specification discloses the same (on page 9, lines 9-15) program 135 (first program)

requires an in-bound connection (e.g. where program 135 is an application server) ... such functionality is provided by a daemon running on computer 105. Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to modify **Aziz et al** to provide a protocol daemon program that does the creating of the first connection to a first program as well as the second connection and coupling the first and second connections thus increasing the functionality of end proxy to allow for other protocols and services as suggested by **Crichton et al** (see column 5, lines 32-45). One skilled in the art would have been lead to make such a modification and recognizes the advantage of using an end proxy that could provide both client and server end proxy functionality as this increase of functionality would allow for more protocols and services as suggested by **Crichton et al** (see column 5, lines 41-45).

As per claim 141, the references as combined above disclose means for transmitting the second connection through a firewall program (see Crichton et al, fig. 4). Claim 141 is rejected on the same rationale as the rejection of claim 140.

As per claim 142, the references as combined above disclose a relay having means for relaying a data stream between the first connection and the second connection (see Aziz et al, column 8, lines 50-65).

As per claim 143, the references as combined above disclose wherein the first security token is one of a password, a network address, and a verification string (see Aziz et al, column 9, lines 3-5).

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As per claim 144, the references as combined above disclose means for terminating the first connection and the second connection (see Aziz et al, column 8, lines 48-49).

Claims 174-178 contain the same limitations as claims 120-124 respectively except for incorporating the claimed method into a computer program. Therefore, 174-178 are rejected on the same rationale as the rejection of claims 120-124.

6. Claims 125-127, 145-147, and 179-181 are rejected under 35 U.S.C. 103(a) as being unpatentable over US Patent 6,643,701 to Aziz et al in view of US Patent 6,104,716 to Crichton et al. as applied to claims 120-124, 140-144, and 174-178 and further in view of US Patent 6,064,671 to Killian.

As per claims 125-127, Crichton et al discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45) that meets the recitation of claim 126. Aziz et al discloses a relay that compares the first security token of the client with stored tokens to perform authentication and in response to a match (i.e. the client's password matches the stored password), coupling the second connection (connection between relay and server) to the first connection which matches the stored security token (see column 12, lines 8-21; column 7, line 65 through column 8, line 1; column 2, lines 11-15). Aziz et al does not explicitly disclose including the second connection

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with one or more corresponding connections in response to no match. Killian in an analogous art teaches providing a plurality of sockets in a socket table or array, wherein each socket normally represents a connection (i.e. has an associated connection) and each socket has an associated network port and network IP address (see column 14, lines 52-64 and column 15, lines 60-67 and column 5, lines 15-22). Killian further discloses means for comparing the network address with the associated network addresses for a match in the socket table (see column 16, lines 6-12 and fig. 3) that meets the recitation of comparing the first security token with the associated security tokens; in response to said comparing if none of the associated security tokens match the first security token (see column 16, lines 17-20), including the socket in the plurality of sockets (see column 16, lines 25-28). Killian also suggests handshaking including use of password (see column 22, lines 3-11). Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to modify the method as combined above to associate the connections with a security token and create new entry in the table in response to no match because it would allow to keep track on the traffic on the network which provides an advantage in load balancing as suggested by Aziz et al (see column 9, lines 31-35 and lines 42-48).

As per claims 145-147, Crichton et al discloses that the functionality of end proxies that meets the recitation of protocol daemon can be increased to allow for other protocols and services, for example one end proxy could provide both client and server end proxy functionality (column 5, lines 41-45) that meets the recitation of claim 146. Aziz et al discloses a relay with means for comparing the first security token of the client with stored tokens to perform

authentication and in response to a match (i.e. the client's password matches the stored password), means for coupling the second connection (connection between relay and server) to the first connection which matches the stored security token (see column 12, lines 8-21; column 7, line 65 through column 8, line 1; column 2, lines 11-15). Aziz et al does not explicitly disclose including the second connection with one or more corresponding connections in response to no match. Killian in an analogous art teaches means for providing a plurality of sockets in a socket table or array, wherein each socket normally represents a connection (i.e. has an associated connection) and each socket has an associated network port and network IP address (see column 14, lines 52-64 and column 15, lines 60-67 and column 5, lines 15-22). Killian further discloses means for comparing the network address with the associated network addresses for a match in the socket table (see column 16, lines 6-12 and fig. 3) that meets the recitation of means for comparing the first security token with the associated security tokens; in response to said comparing if none of the associated security tokens match the first security token (see column 16, lines 17-20), means for including the socket in the plurality of sockets (see column 16, lines 25-28). Killian also suggests handshaking including use of password (see column 22, lines 3-11). Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to modify the method as combined above to associate the connections with a security token and create new entry in the table in response to no match because it would allow to keep track on the traffic on the network which provides an advantage in load balancing as suggested by Aziz et al (see column 9, lines 31-35 and lines 42-48).

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Claims 179-181 contains the same limitations as claims 125-127 respectively except for incorporating the claimed method into a computer program. Therefore, 179-181 are rejected on the same rationale as the rejection of claims 125-127.

Conclusion

7. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Carl Colin whose telephone number is 571-272-3862. The examiner can normally be reached on Monday through Thursday, 8:00-6:30 PM.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Nasser G. Moazzami can be reached on 571-272-4195. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

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Carl Colin Patent Examiner February 16, 2007 NASSER MOAZZAMI SUPERVISORY PATENT EXAMINER TECHNOLOGY CENTER 2100

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